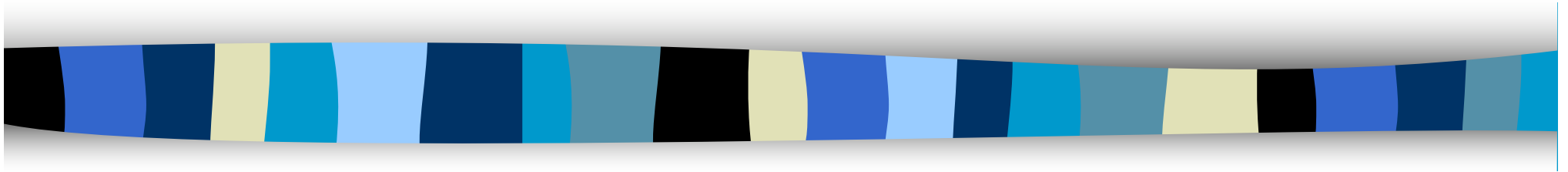


Similarity Search for Adaptive Ellipsoid Queries Using Spatial Transformation



Yasushi Sakurai (*NTT Cyber Space Laboratories*)

Masatoshi Yoshikawa (*Nara Institute of Science and Technology*)

Ryoji Kataoka (*NTT Cyber Space Laboratories*)

Shunsuke Uemura (*Nara Institute of Science and Technology*)



Outline

- Introduction
- STT (spatial transformation technique)
 - Definition of spatial transformation
 - Spatial transformation of rectangles
 - Search algorithm
- MSTT (multiple STT)
 - Index structure construction
 - Query processing
 - Dissimilarity of matrices
- Performance test
- Conclusion

Introduction

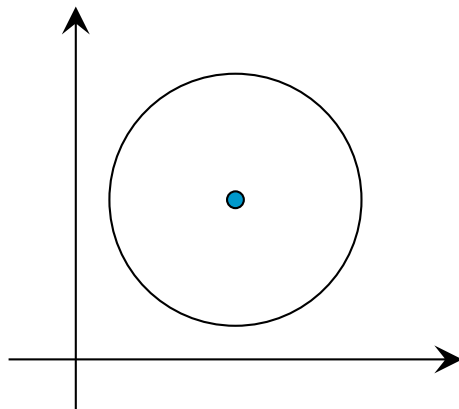
■ Ellipsoid query

- Search processing is performed by using quadratic form distance functions

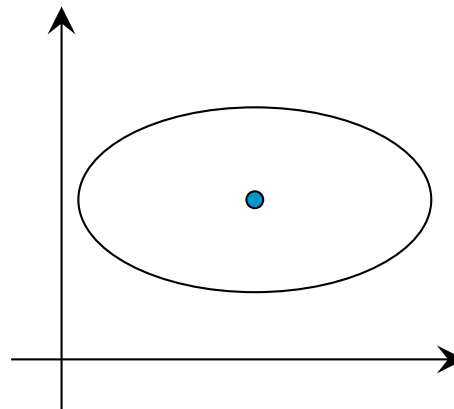
- Distance of p and q for a query matrix M :

$$d_M^2(p, q) = (p - q) \cdot M \cdot (p - q)^t$$

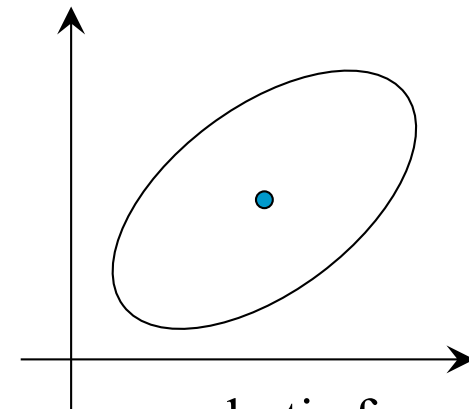
- represents correlations between dimensions



Euclidean
circles for isosurfaces



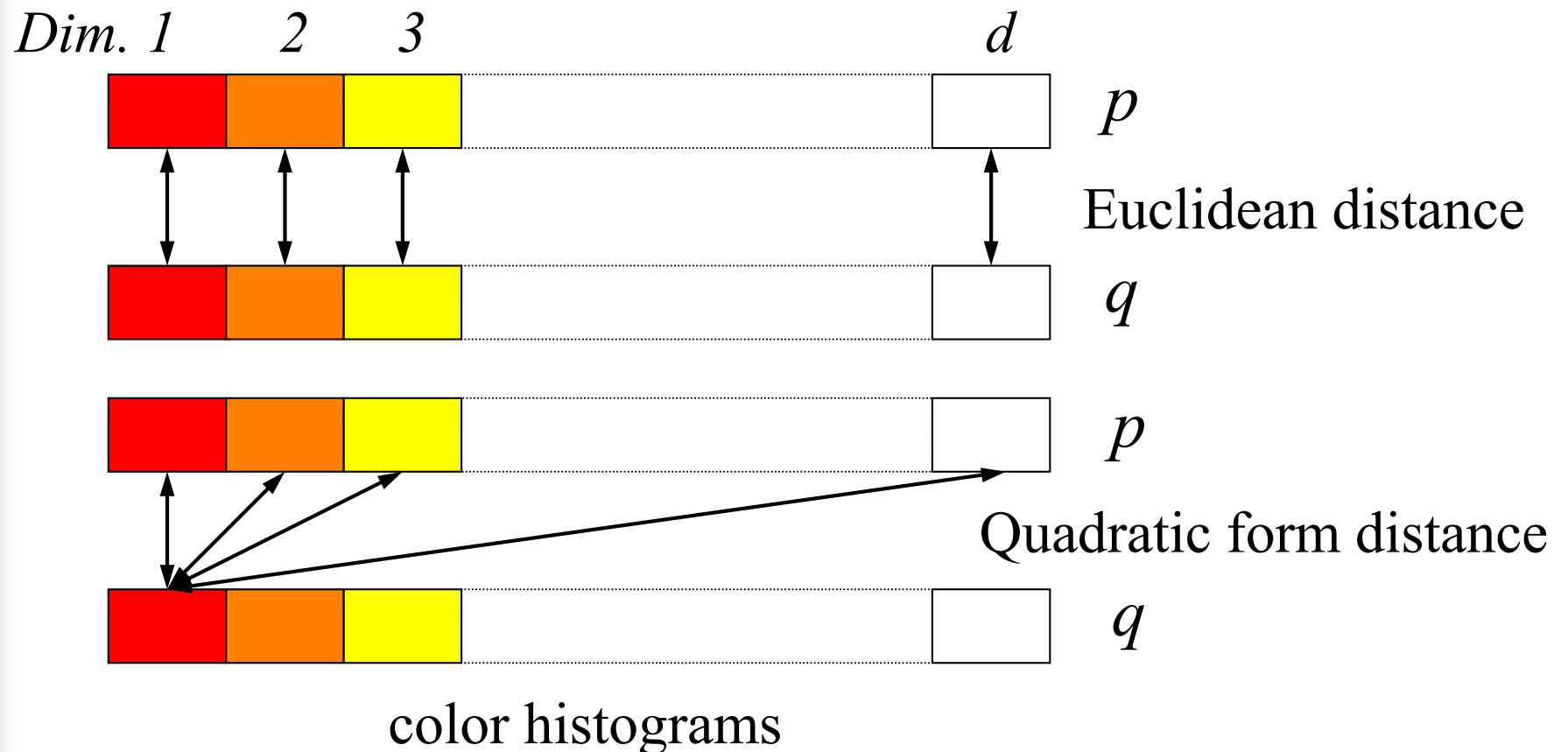
weighted Euclidean
iso-oriented ellipsoids



quadratic form
Ellipsoids
(Not necessarily aligned
to the coordinate axis)

Introduction

- An application of a quadratic form distance function
 - represent the similarity between colors i and j



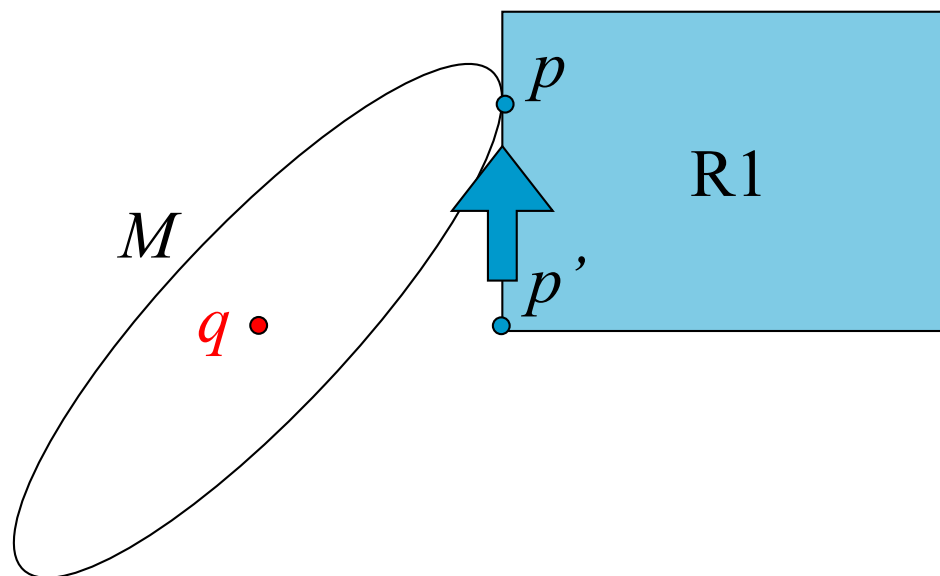


Introduction

- Spatial indices
 - e.g. R-tree family (R*-tree, X-tree, SR-tree, A-tree)
 - Based on the Euclidean distance function
 - ➡ Cannot be applied to ellipsoid queries
- Efficient search methods for **user-adaptive ellipsoid queries**
 - Query matrix M is variable

Related Work : Seidl and Kriegel, VLDB97

- Search method based on the steepest descent method
 - Works on spatial indices of R-tree family
 - Calculates the exact distance of a query point and an MBR in an index structure
 - ...but requires high CPU cost which exceeds disk access cost



Moves p' toward p iteratively

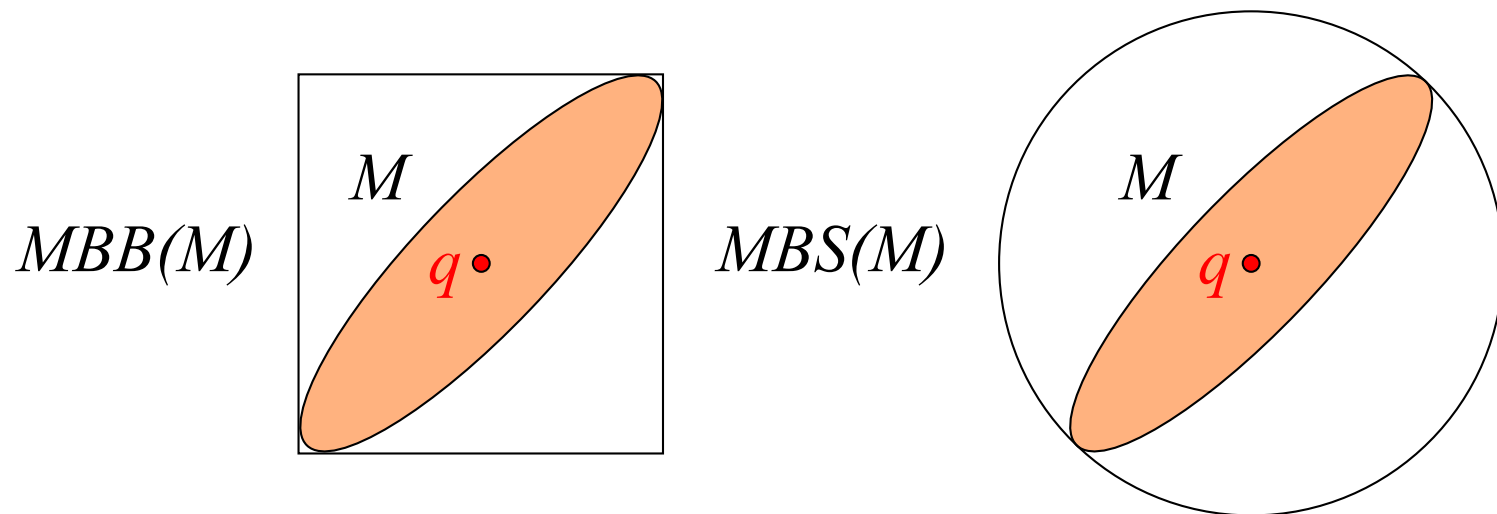
CPU time $O(\omega d^2)$
 ω ...number of iterations
 d ...dimensionality

Related Work : Ankerst et al., VLDB98

- Technique that uses the MBB and MBS distance functions to reduce CPU time
 - MBB and MBS distance functions

$$d_{MBB(M)}^2(p, q) = \max_{i=1}^d \left((p_i - q_i)^2 / (M^{-1})_{ii} \right)$$

$$d_{MBS(M)}^2(p, q) = \lambda_{M_{\min}}^2 \cdot (p - q)^2$$





Related Work : Ankerst et al., VLDB98

- Approximation technique by using the MBB and MBS distance functions
 - approximation distance : uses either MBB or MBS distance for better approximation quality
 - Calculates the exact distances only if data objects or MBRs cannot be filtered by their approximation distances
 - Saves CPU time by reducing the number of exact distance calculations
 - ...but cannot reduce the number of exact distance calculations if its approximation quality is low



Our Contributions

- STT (Spatial Transformation Technique)
 - Ellipsoid queries incur a high CPU cost
 - The efficiency depends on approximation quality
 - STT efficiently processes ellipsoid queries because of high approximation quality
- MSTT (Multiple Spatial Transformation Technique)
 - Does not use only the Euclidean distance function to make index structures
 - Ellipsoid queries give various distance functions
 - In MSTT, various index structures are created; the search algorithm utilizes a structure well suited to a query matrix

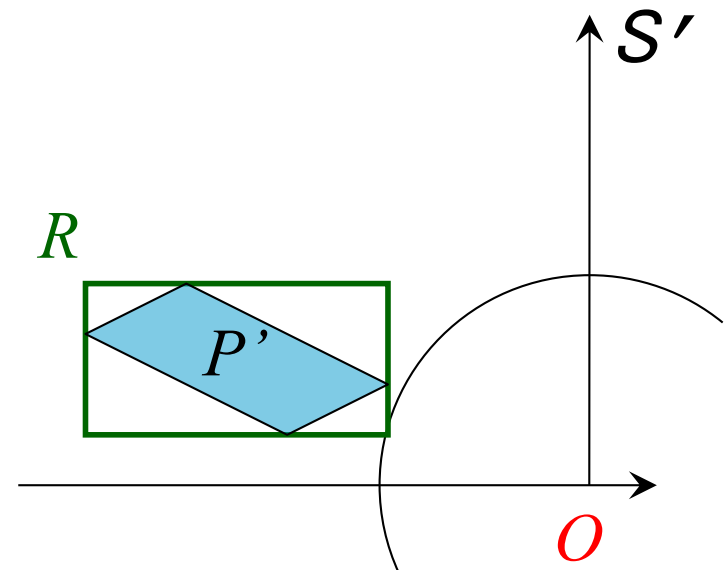
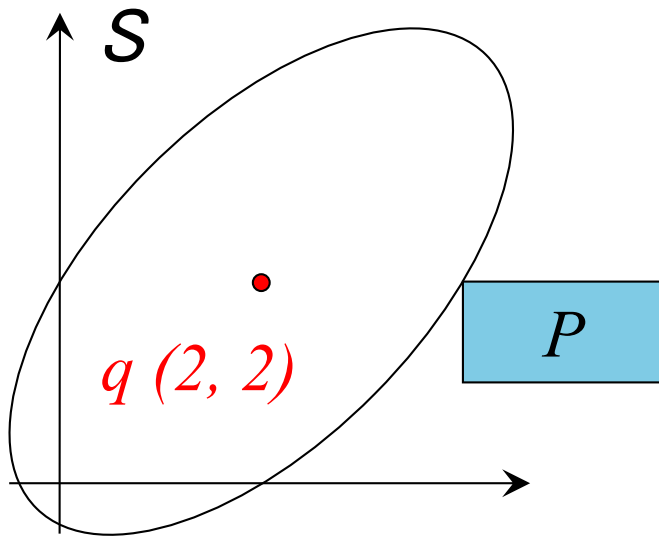


Outline

- Introduction
- **STT (spatial transformation technique)**
 - Definition of spatial transformation
 - Spatial transformation of rectangles
 - Search algorithm
- MSTT (multiple STT)
 - Index structure construction
 - Query processing
 - Dissimilarity of matrices
- Performance test
- Conclusion

Spatial Transformation Technique (STT)

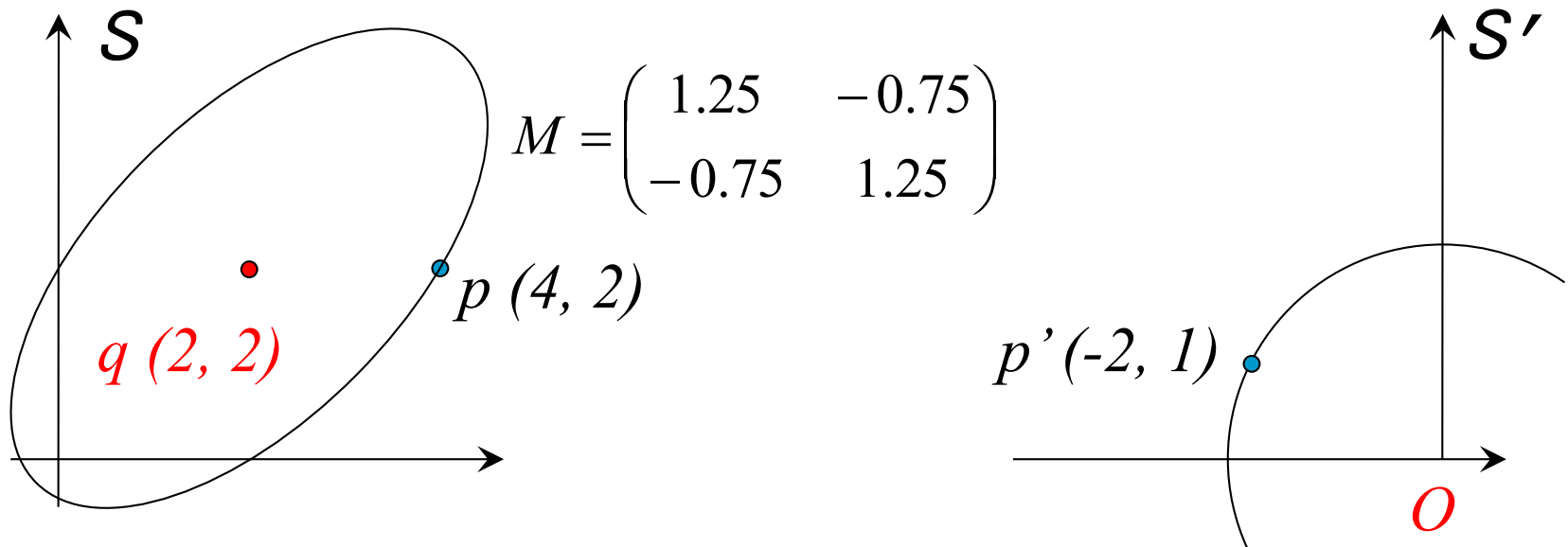
- High approximation quality
 - STT consumes less CPU time
- Spatial transformation
 - MBRs in a quadratic form distance space are transformed into rectangles in the Euclidean distance space



Spatial Transformation

■ Definition of spatial transformation

- p : a point in the quadratic form distance space S
- p' : a point in the Euclidean distance space S'
- The distance between q and p in S is equal to the distance between p' and O in S'
- Spatial transformation of p into p'



Spatial Transformation

■ Definition of spatial transformation

- $d_M^2(p, q)$: the distance of p and q in S

$$d_M^2(p, q) = (p - q) \cdot M \cdot (p - q)^t$$

- E_M : the eigenvector of M , Λ_M : the eigenvalues of M

$$M = E_M \cdot \Lambda_M \cdot E_M^t$$

$$d_M^2(p, q) = (p - q) \cdot E_M \cdot \Lambda_M \cdot E_M^t \cdot (p - q)^t$$

- Spatial transformation of p into p'

$$d_M^2(p, q) = p' \cdot p'^t = d^2(p', O)$$

$$p' = (p - q) \cdot A_M$$

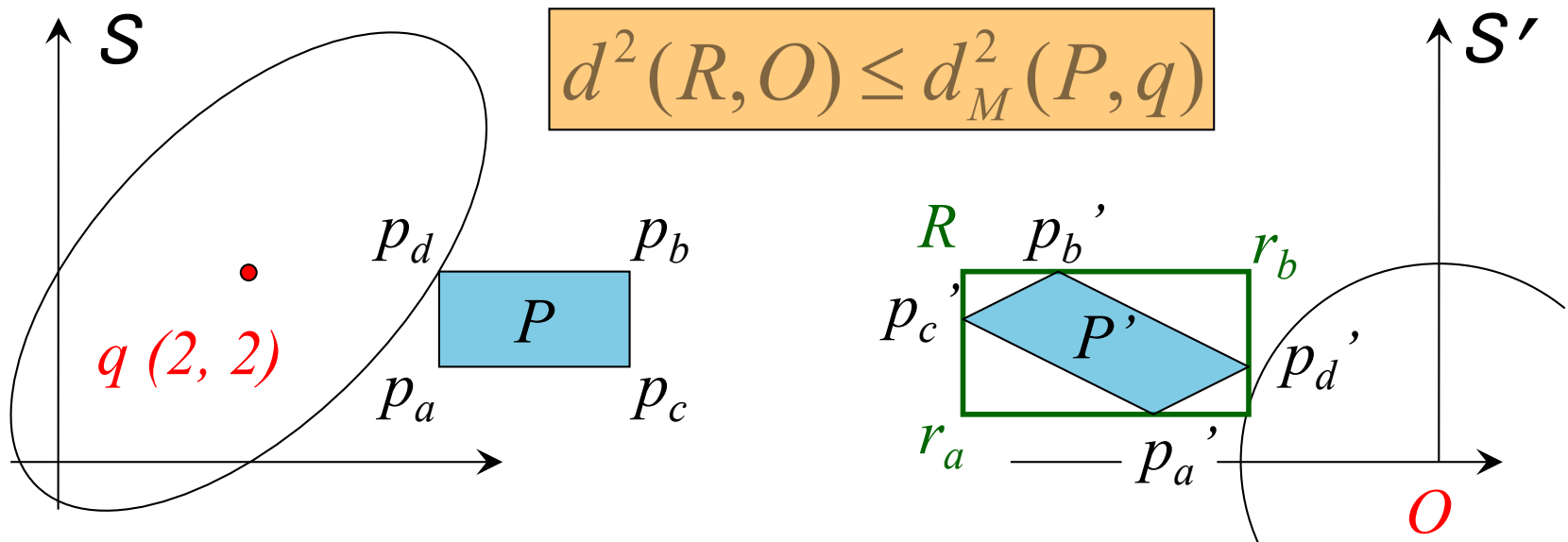
$$A_M = E_M \cdot \Lambda_M^{1/2}$$

Approximation Rectangles

1. P in S is transformed into P' in S'

The calculation of distance between the origin and polygons in high-dimensional spaces incurs a high CPU cost

2. P' is approximated by R
 3. $d^2(R, O)$ is used instead of $d_M^2(P, q)$
- } low CPU cost



Approximation Rectangles

1. Calculates $p'_a = (p_a - q) \cdot A_M$

p_a : lower endpoint of the major diagonal of P

2. Creates the two matrices from the components

a_{ij} of A_M

$$\phi_{ij} = \begin{cases} a_{ij} & (a_{ij} < 0) \\ 0 & (otherwise), \end{cases} \quad \psi_{ij} = \begin{cases} a_{ij} & (a_{ij} > 0) \\ 0 & (otherwise) \end{cases}$$

3. Calculates the approximation rectangle R of P'

$$R = (r_a, r_b),$$

$$r_{a_j} = p'_{a_j} + \sum_{i=1}^d l_i \cdot \phi_{ij}, \quad r_{b_j} = p'_{a_j} + \sum_{i=1}^d l_i \cdot \psi_{ij}$$

l_i : the edge length of P for the i -th dimension

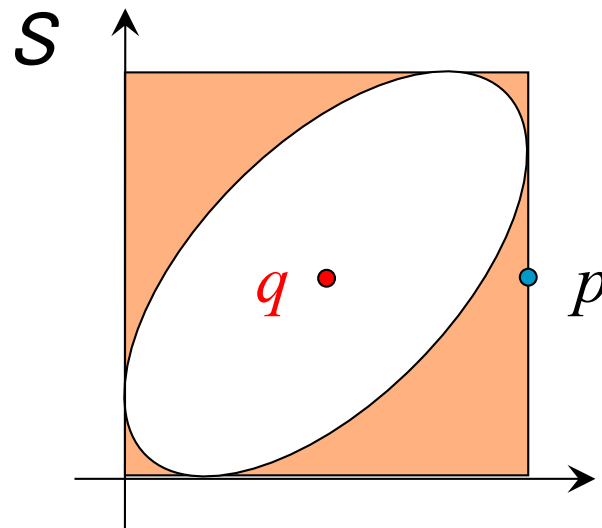
4. R can be used for search since R totally contains P' , that is $d^2(R, O) \leq d_M^2(P, q)$

Search Algorithm

1. Calculates the transformation matrix of M
2. Searches for similarity objects by using an index

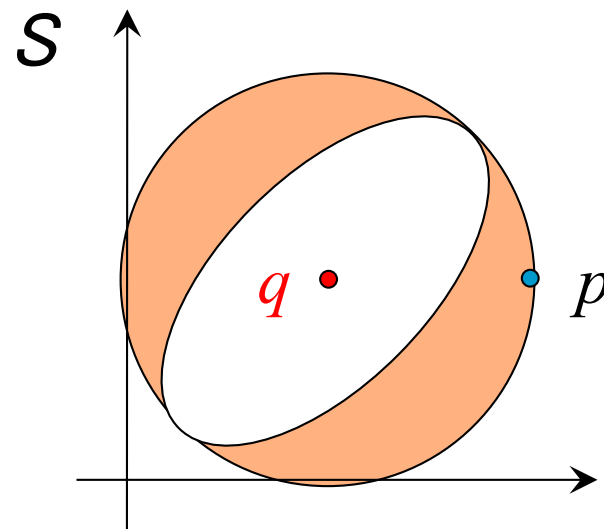
[Data nodes]

– Calculates $d_{MBB-MBS(M)}(p, q)$



Search Algorithm

1. Calculates the transformation matrix of M
2. Searches for similarity objects by using an index
[Data nodes]
 - Calculates $d_{MBB-MBS(M)}(p, q)$

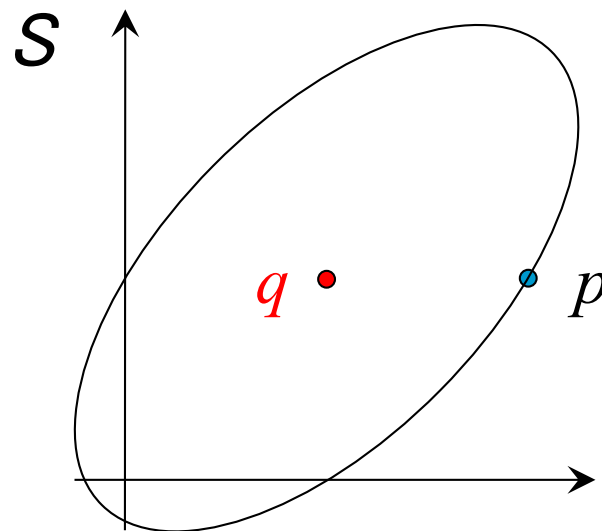


Search Algorithm

1. Calculates the transformation matrix of M
2. Searches for similarity objects by using an index

[Data nodes]

- Calculates $d_{MBB-MBS(M)}(p, q)$
- Calculates $d_M(P, q)$ if $d_{MBB-MBS(M)}(p, q) \leq d_{(M)}(k\text{-}NN, q)$

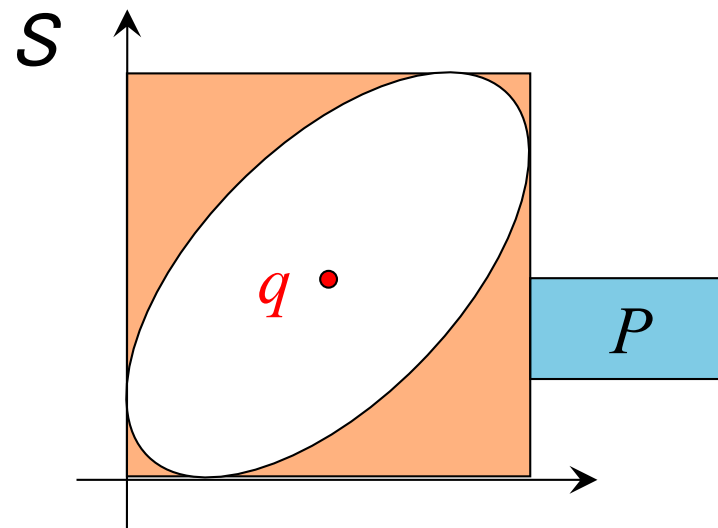


Search Algorithm

1. Calculates the transformation matrix of M
2. Searches for similarity objects by using an index

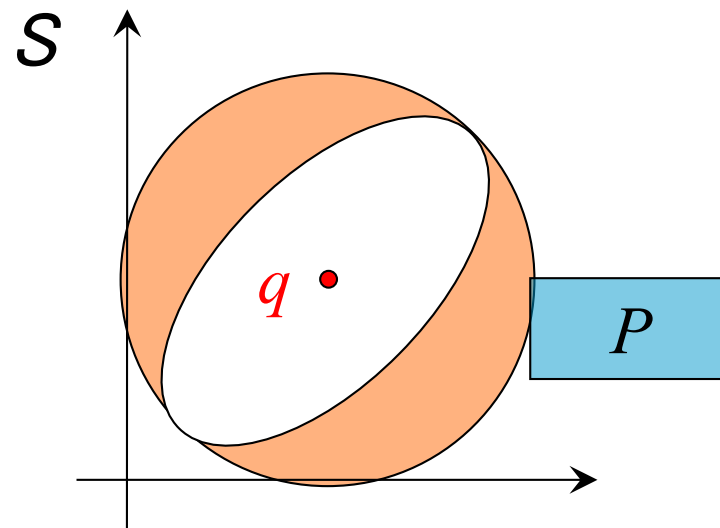
[Directory nodes]

– Calculates $d_{MBB-MBS(M)}(P, q)$



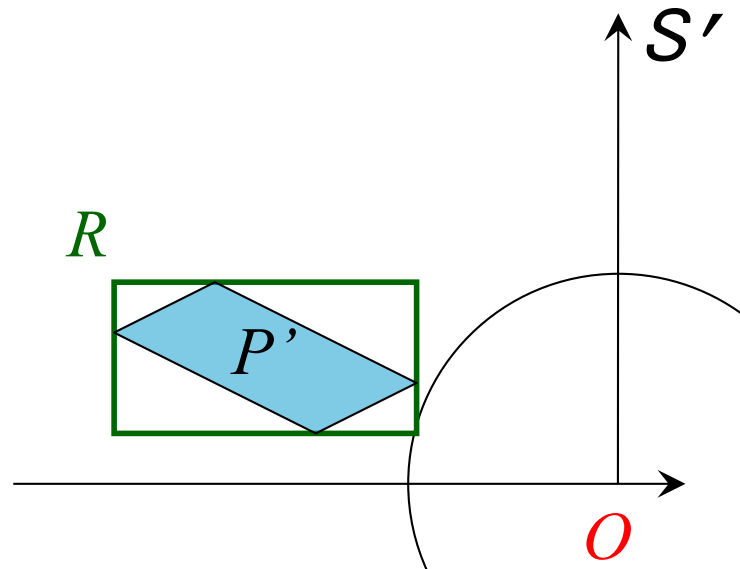
Search Algorithm

1. Calculates the transformation matrix of M
2. Searches for similarity objects by using an index
[Directory nodes]
 - Calculates $d_{MBB-MBS(M)}(P, q)$



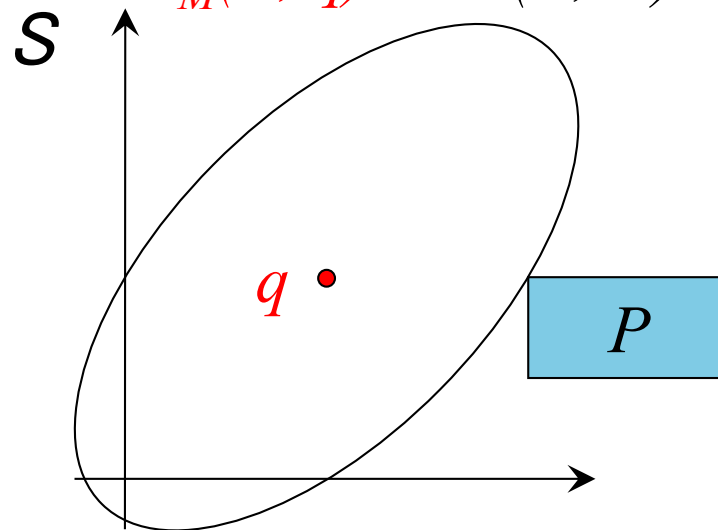
Search Algorithm

1. Calculates the transformation matrix of M
2. Searches for similarity objects by using an index
[Directory nodes]
 - Calculates $d_{MBB-MBS(M)}(P, q)$
 - Calculates $d(R, O)$ if $d_{MBB-MBS(M)}(P, q) \leq d_{(M)}(k\text{-}NN, q)$



Search Algorithm

1. Calculates the transformation matrix of M
2. Searches for similarity objects by using an index
[Directory nodes]
 - Calculates $d_{MBB-MBS(M)}(P, q)$
 - Calculates $d(R, O)$ if $d_{MBB-MBS(M)}(P, q) \leq d_{(M)}(k\text{-NN}, q)$
 - Calculates $d_M(P, q)$ if $d(R, O) \leq d_{(M)}(k\text{-NN}, q)$





Outline

- Introduction
- STT (spatial transformation technique)
 - Definition of spatial transformation
 - Spatial transformation of rectangles
 - Search algorithm
- **MSTT (multiple STT)**
 - Index structure construction
 - Query processing
 - Dissimilarity of matrices
- Performance test
- Conclusion



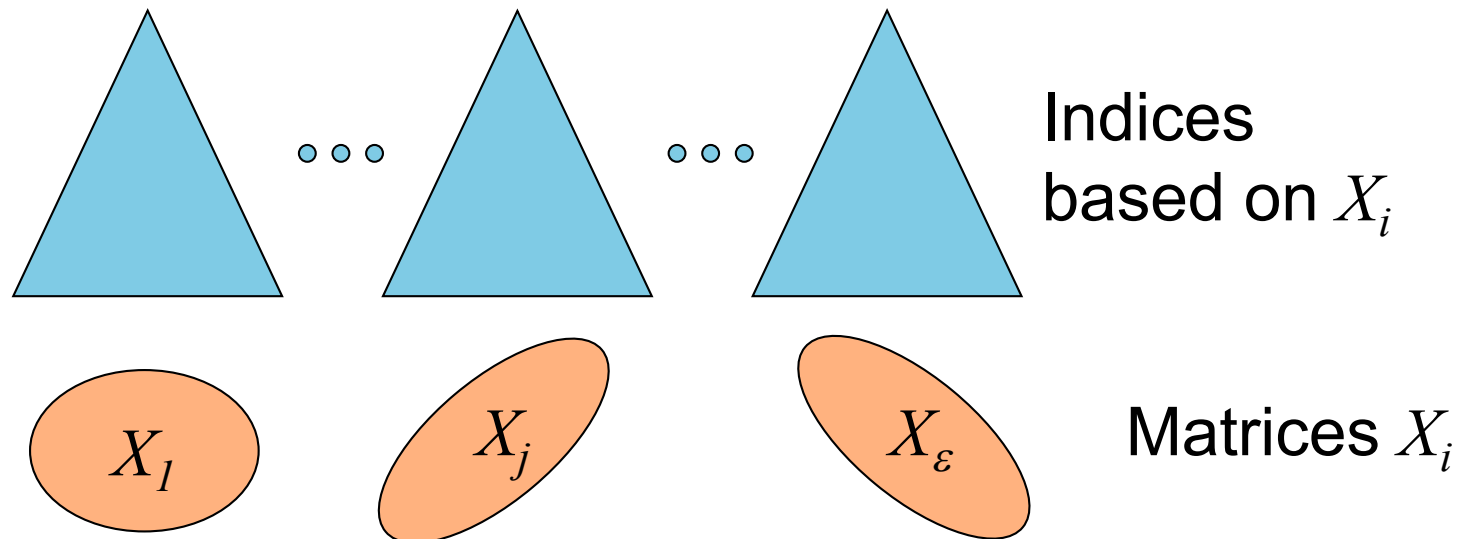
Multiple Spatial Transformation Technique (MSTT)

- Node access problem
 - If a query matrix is NOT similar to the unit matrix, it causes a large number of node accesses
 - Index structures are constructed by the Euclidean distance function
- Constructs various index structures by using quadratic form distance functions
 - Chooses a structure that gives sufficient search performance in query processing
 - Reduces both CPU time and number of page accesses for ellipsoid queries

Basic Idea

■ Similarity of matrices

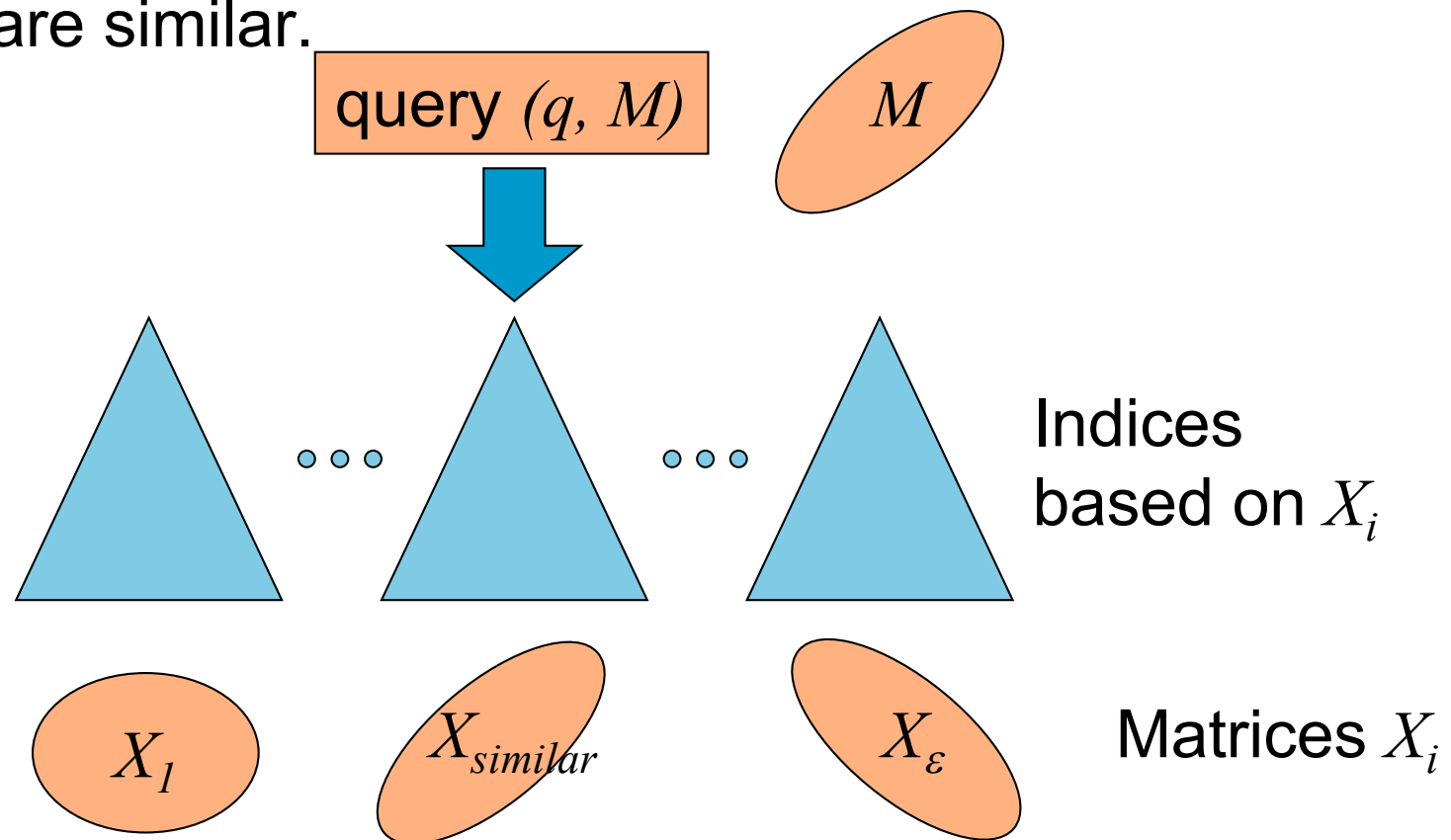
- High search performance can be expected when the query matrix and the matrix of selected index are similar.



Basic Idea

■ Similarity of matrices

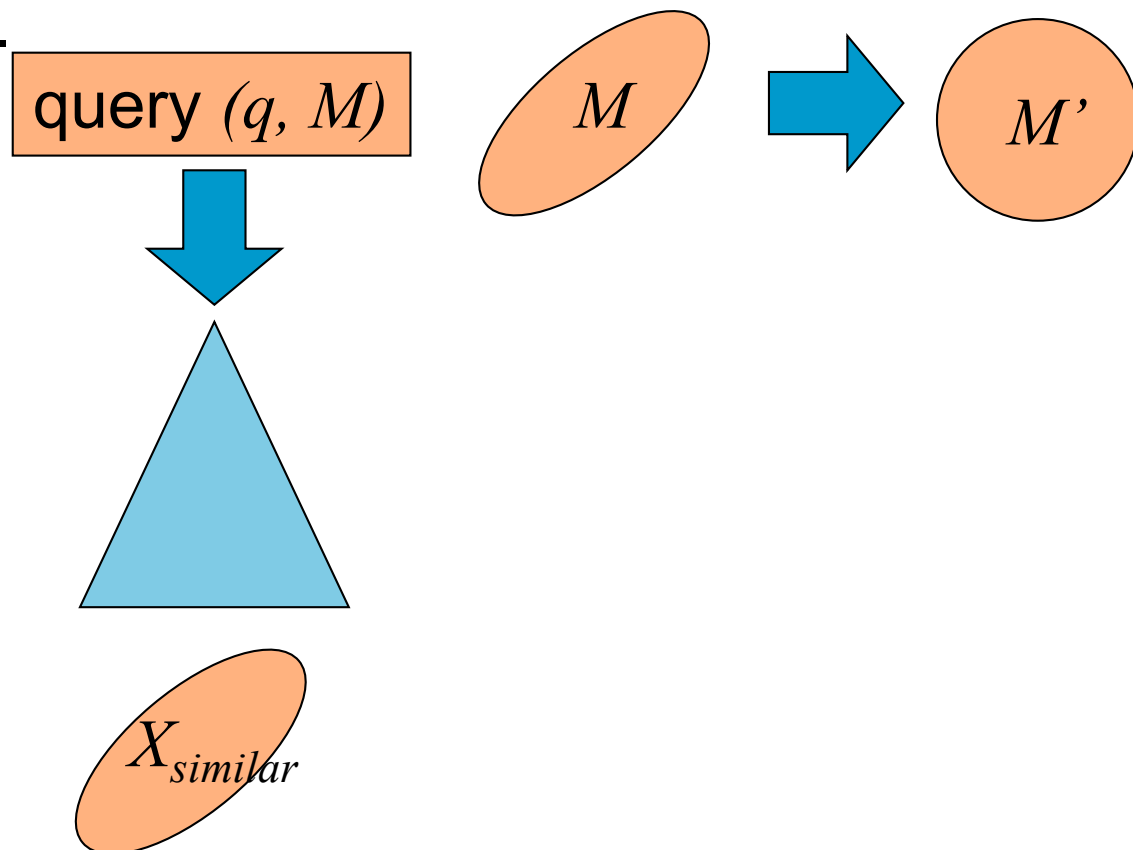
- High search performance can be expected when the query matrix and the matrix of selected index are similar.



Basic Idea

■ Similarity of matrices

- High search performance can be expected when the query matrix and the matrix of selected index are similar.





Indexing and Retrieval Mechanism

■ Index structure construction

- C : the matrix for constructing the index I_C
- Transformation matrix $A_C = E_C \cdot \Lambda_C^{1/2}$
- All data points in a data set are transformed
$$p' = p \cdot A_C$$
- I_C is constructed using transformed data points

Indexing and Retrieval Mechanism

■ Query processing

1. Calculates the transformed query point $q' = q \cdot A_C$

2. Calculates the query matrix

$$M' = A_C^{-1} \cdot M \cdot (A_C^{-1})^t$$

3. Performs search processing by using I_C, M', q'

■ The query of M can be processed by using I_C

$$\begin{aligned} d_M^2(p, q) &= (p - q) \cdot M \cdot (p - q)^t \\ &= (p' - q') \cdot A_C^{-1} \cdot M \cdot (A_C^{-1})^t \cdot (p' - q')^t \\ &= (p' - q') \cdot M' \cdot (p' - q')^t \end{aligned}$$



Similarity of Matrices

■ Flatness of a query matrix

- The variance σ_M^2 of the eigenvalues of M is called the flatness of M :

$$\sigma_M^2 = \sum_{i=1}^d (\lambda_{M_i} - \bar{\lambda}_M)^2, \quad \bar{\lambda}_M = \sum_{j=0}^d \lambda_{M_j} / d$$

λ_{M_i} : the i -th dimensional eigenvalue

$\bar{\lambda}_M$: the average of the eigenvalues of M

- The flatness of the unit matrix is 0 (search of the Euclidean space).



Similarity of Matrices

■ Dissimilarity of M and I_C

- MSTT employs $\sigma_{M'}^2$ as the measure of the dissimilarity between M and I_C
- $\sigma_{M'}^2$: the flatness of M'
$$M' = A_C^{-1} \cdot M \cdot (A_C^{-1})^t$$
- The effectiveness of I_c relative to M improves as $\sigma_{M'}^2$ decreases



Outline

- Introduction
- STT (spatial transformation technique)
 - Definition of spatial transformation
 - Spatial transformation of rectangles
 - Search algorithm
- MSTT (multiple STT)
 - Index structure construction
 - Query processing
 - Dissimilarity of matrices
- Performance test
- Conclusion



Performance Test

- Data sets: real data set (rgb histogram of images)
- Data size: 100,000
- Dimensionality: 8 and 27
- Page size: 8 KB
- 20-nearest neighbor queries
- Evaluation is based on the average for 100 query points
- Index structure : A-tree (Sakurai et al., VLDB2000)
- CPU: SUN UltraSPARC-II 450MHz



Performance Test

■ Query matrices for experiments

- [HSE⁺95] : the components of M

$$m_{ij} = \exp(-\alpha(d_w(c_i, c_j)/d_{\max})^2)$$

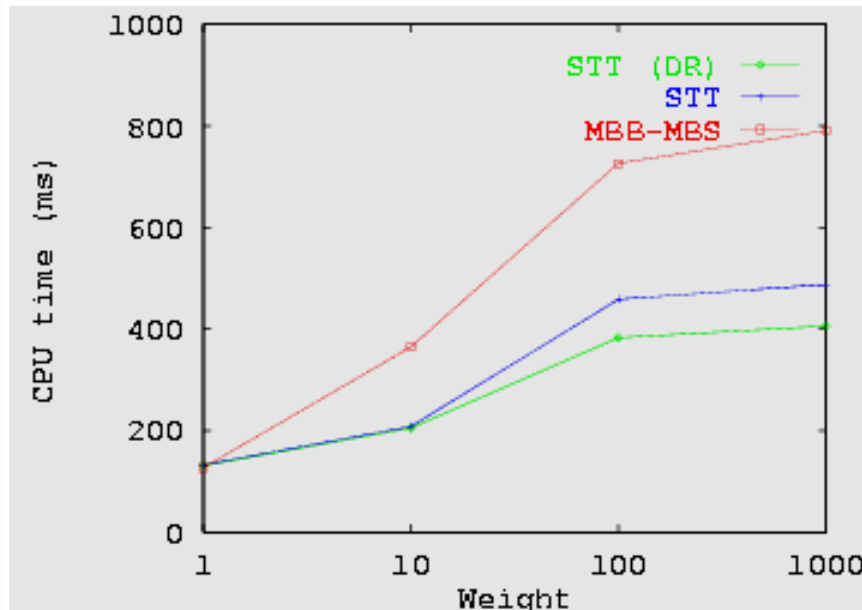
α : positive constant,

$d_w(c_i, c_j)$: the weighted Euclidean distance
between the color c_i and c_j ,

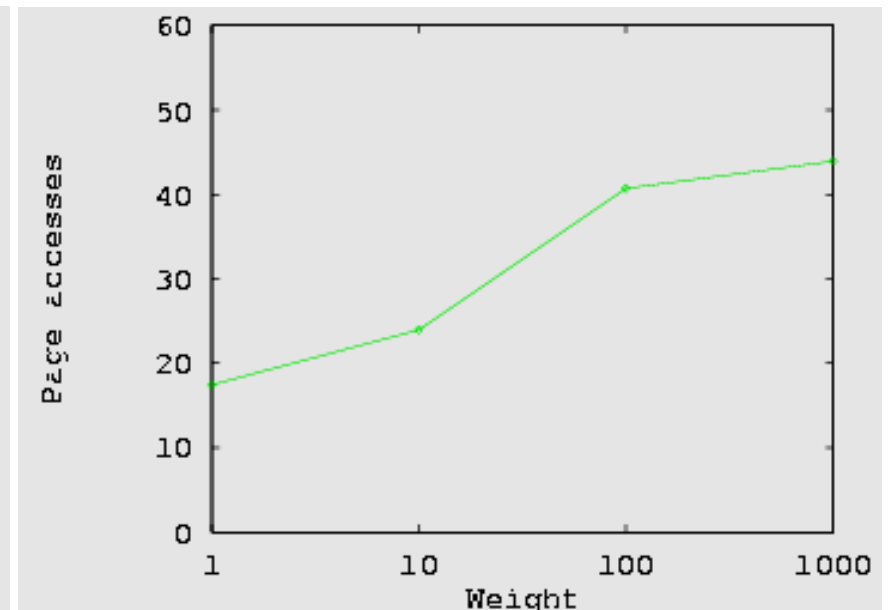
$w=(w_r, w_g, w_b)$: the weightings of the red, green
and blue components in RGB color space

- $\alpha=10, w_g=w_b=1$
- w_r was varied from 1 to 1,000
- The flatness of M increases as w_r becomes large

Performance of STT



CPU time ($d = 8$)

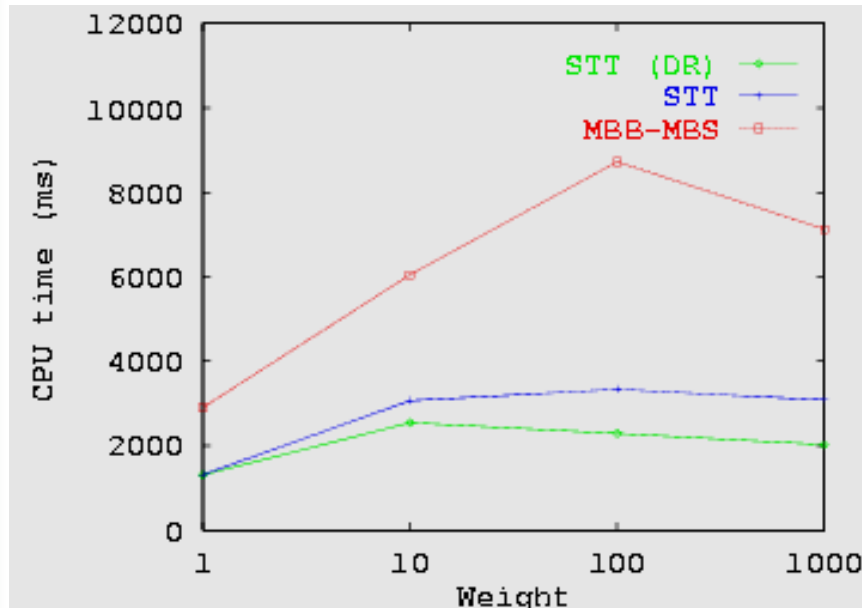


Number of page accesses ($d = 8$)

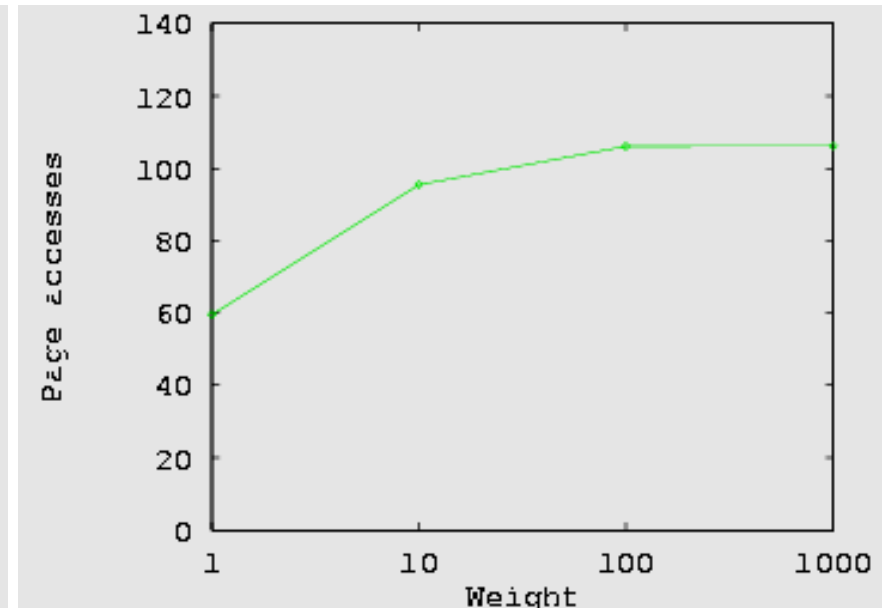
■ Comparison of STT and MBB-MBS (8D)

- Both methods require the same number of page accesses since they utilize exact distance functions
- Low CPU cost : STT increases approximation quality, and reduces the number of exact calculations
- The effectiveness of STT increases with matrix flatness

Performance of STT



CPU time ($d = 27$)

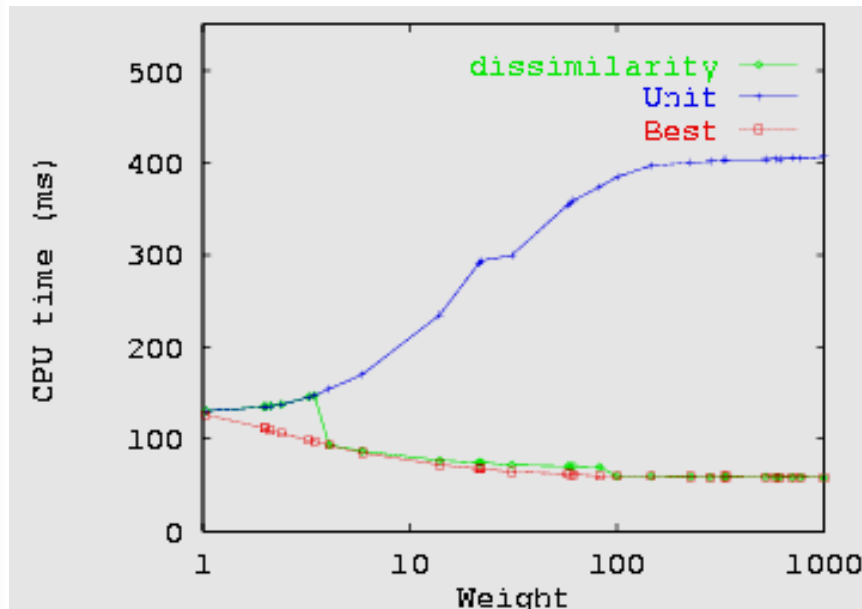


Number of page accesses ($d = 27$)

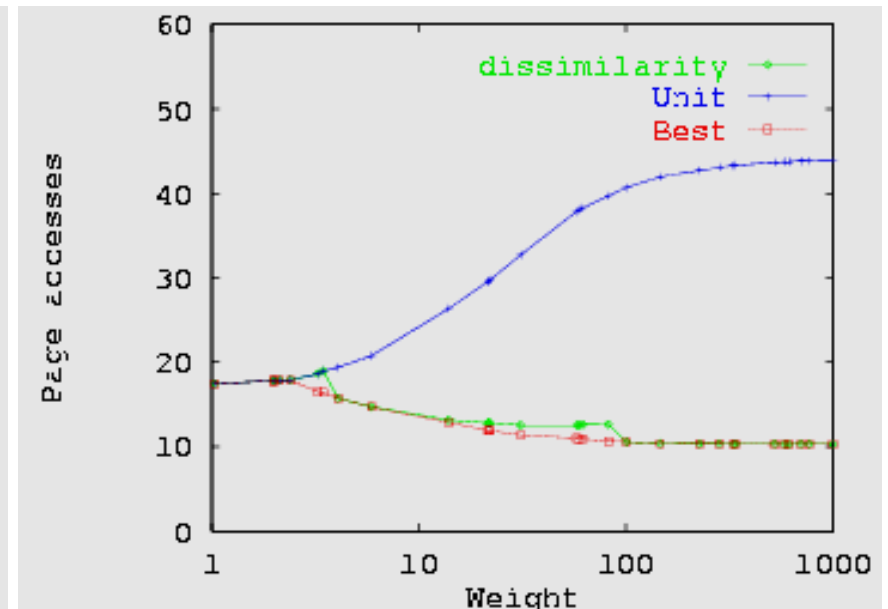
■ Comparison of STT and MBB-MBS (27D)

- The effectiveness of STT increases as either dimensionality or matrix flatness grows
- STT achieves a 74% reduction in CPU cost for high dimensionality and matrix flatness

Performance of MSTT



CPU time ($d = 8$)



Number of page accesses ($d = 8$)

■ Three structures

- structure constructed by the unit matrix (**Unit**)
- structure constructed by the matrix $w_r=10$
- structure constructed by the matrix $w_r=1000$

■ Performance of MSTT

- Dissimilarity : the cost of search using a structure chosen by the dissimilarity function
- Dissimilarity is not optimal, but provides good performance

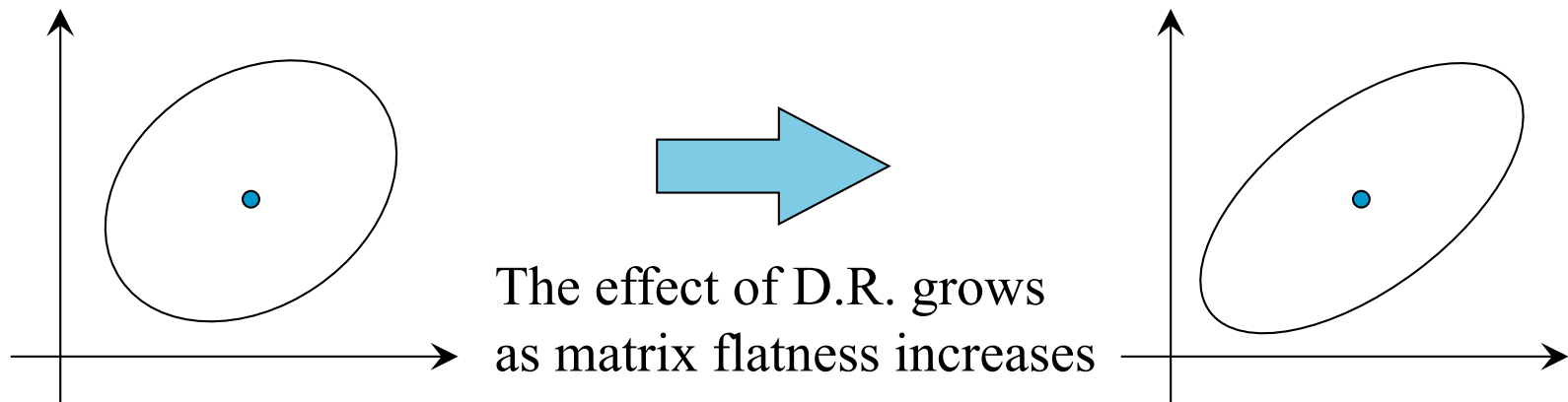


Conclusions

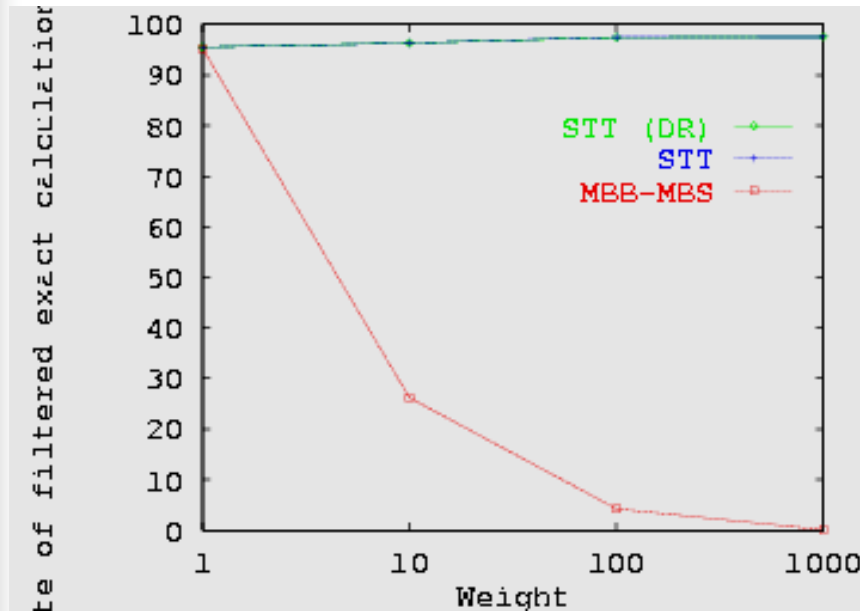
- Search methods for user-adaptive ellipsoid queries
- STT (Spatial Transformation Technique)
 - Spatial transformation : MBRs in the quadratic form distance space are transformed into rectangles in the Euclidean distance space
 - STT performs ellipsoid queries efficiently even when dimensionality or matrix flatness is high
- MSTT (Multiple Spatial Transformation Technique)
 - MSTT creates various index structures; the search algorithm utilizes a structure well suited to a query matrix
 - MSTT reduces both CPU time and the number of page accesses

Dimensionality Reduction

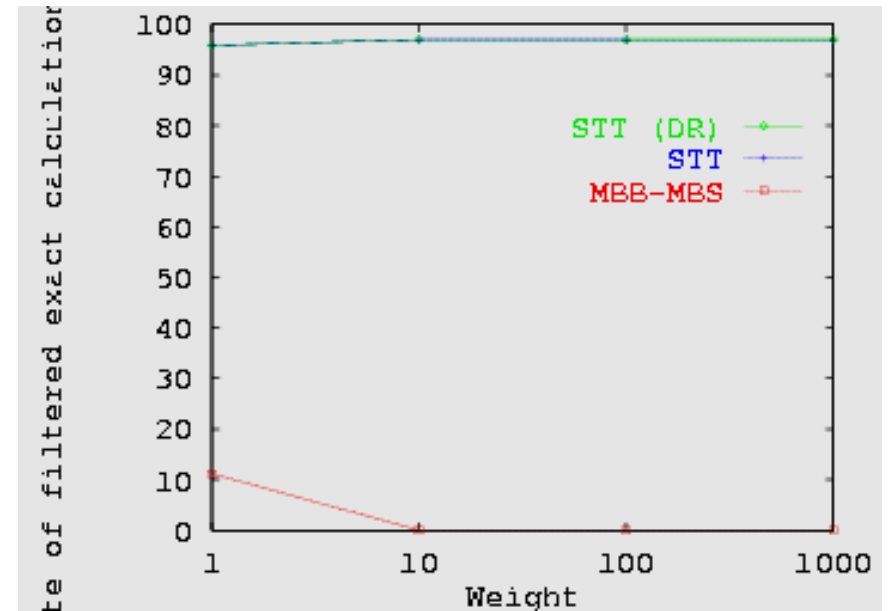
- Eigenvalues of a query matrix
 - Dimensions corresponding to small eigenvalues contribute less to approximation quality
 - These dimensions are eliminated to save on CPU costs
 - Calculation time for the spatial transformation of rectangles is reduced to n/d ($n \geq d$)
 n : the number of dimensions used



Performance of STT (2)



$d = 8$

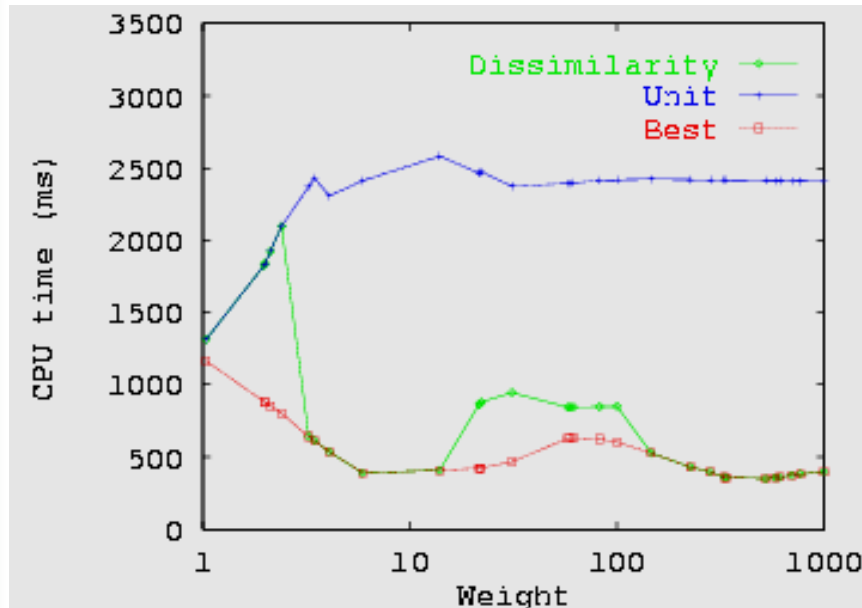


$d = 27$

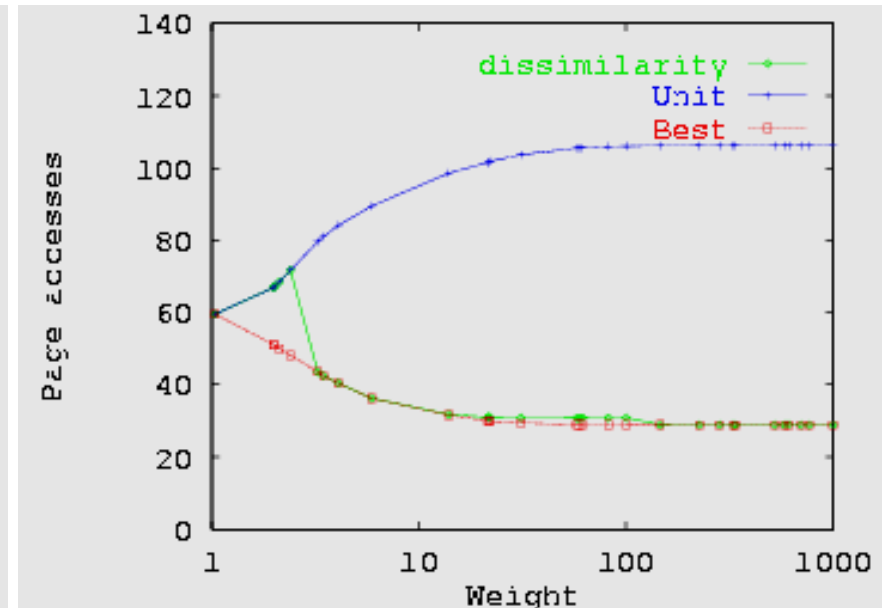
Rate of filtered exact calculations

- Percentage of filtered exact distance calculations
 - The efficiency of MBB-MBS decreases as matrix flatness grows
 - STT effectively filters exact distance calculations for all queries

Performance of MSTT



CPU time ($d = 27$)



Number of page accesses ($d = 27$)

■ Low search cost

- Compared with the structure by the Euclidean distance function, MSTT reduces both CPU time and the number of page accesses
- MSTT constructs various structures
- Dissimilarity function chooses structures well suited to the query matrix.